

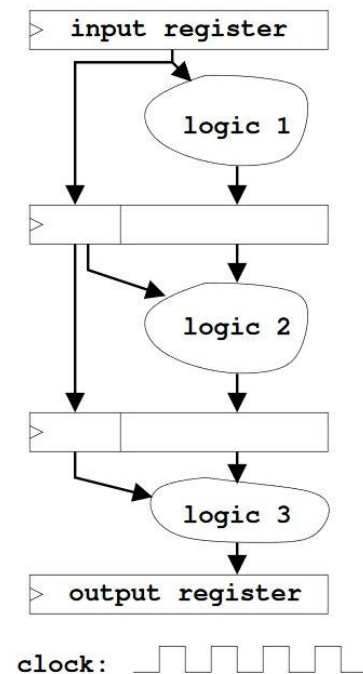
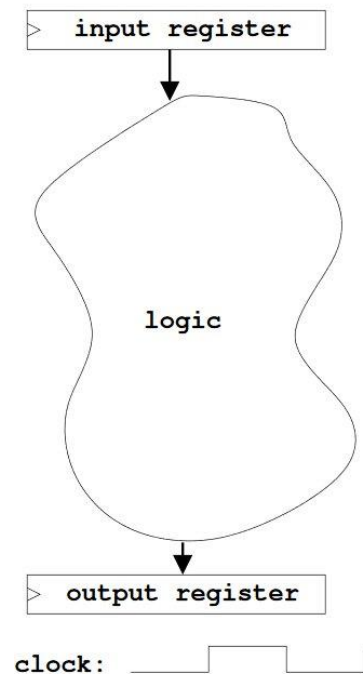
Tutorial 12

December 20, 2016

Speed up the β -machine with pipelines

- The concept of pipeline is useful to increase the rate at which a deep combinatorial circuit can process the data

$$MIPS = \frac{Freq}{CPI}$$



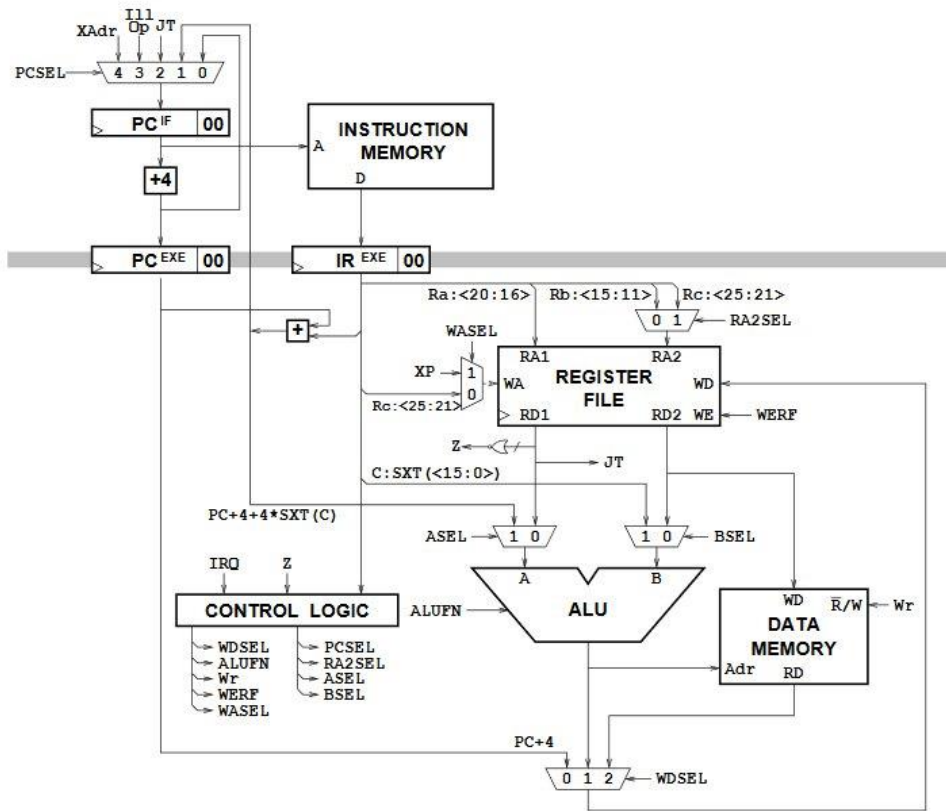
- **Terminology:**
 - **MIPS:** millions of instruction per second
 - **Freq:** clock frequency
 - **CPI:** clock (cycle) per instruction
- To increase performances we need either to **increase the frequency** or to **decrease the CPI**
- We cannot increase the frequency because we are blocked by the stabilization time of the circuit
- We can decrease the CPI \Rightarrow pipelining

Pipelining the β -machine

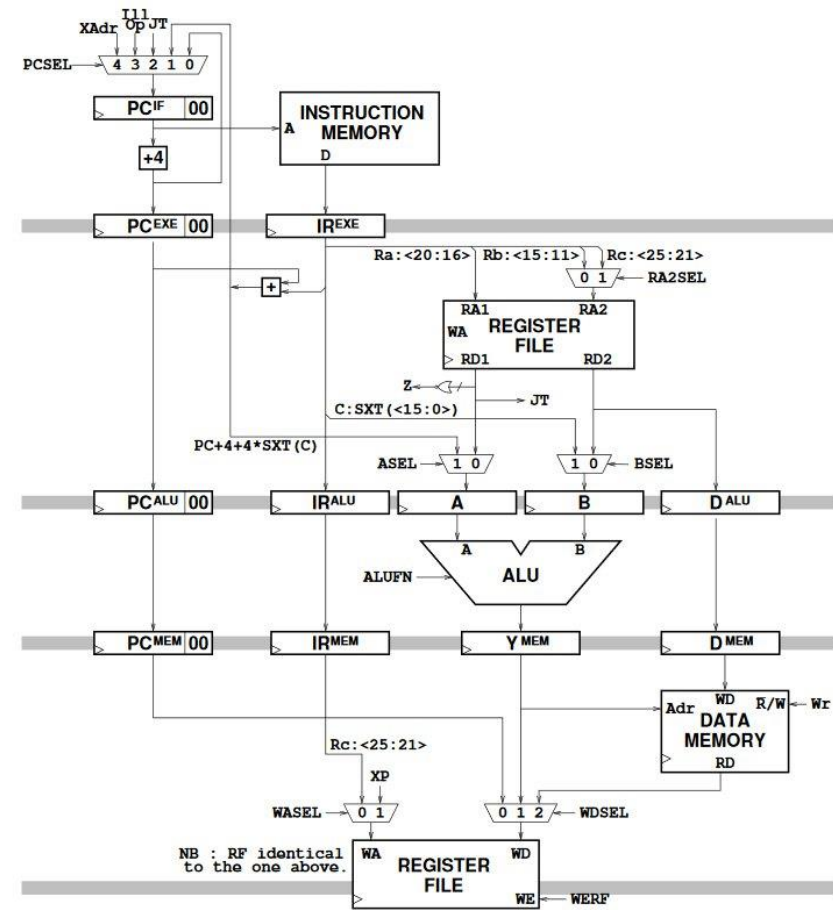
- We can decompose the combinatorial logic into five identifiable stages:
 - **Instruction fetch (IF)**: fetch the instruction from memory
 - **Register fetch (RF)**: fetch the arguments from the registers
 - **ALU**: computation of the ALU
 - **Memory (MEM)**: reading data from memory
 - **Write back (WB)**: writing back the results to the registers
- During this tutorial, we will study simplified versions with two stages and four stages (see next slide)

Pipelining the β -machine

A β machine with a 2-stage pipeline



2 stages: {IF}, {RF, ALU, MEM, WB}



4 stages: {IF}, {RF}, {ALU}, {MEM, WB}

Pipelining comes at a cost...

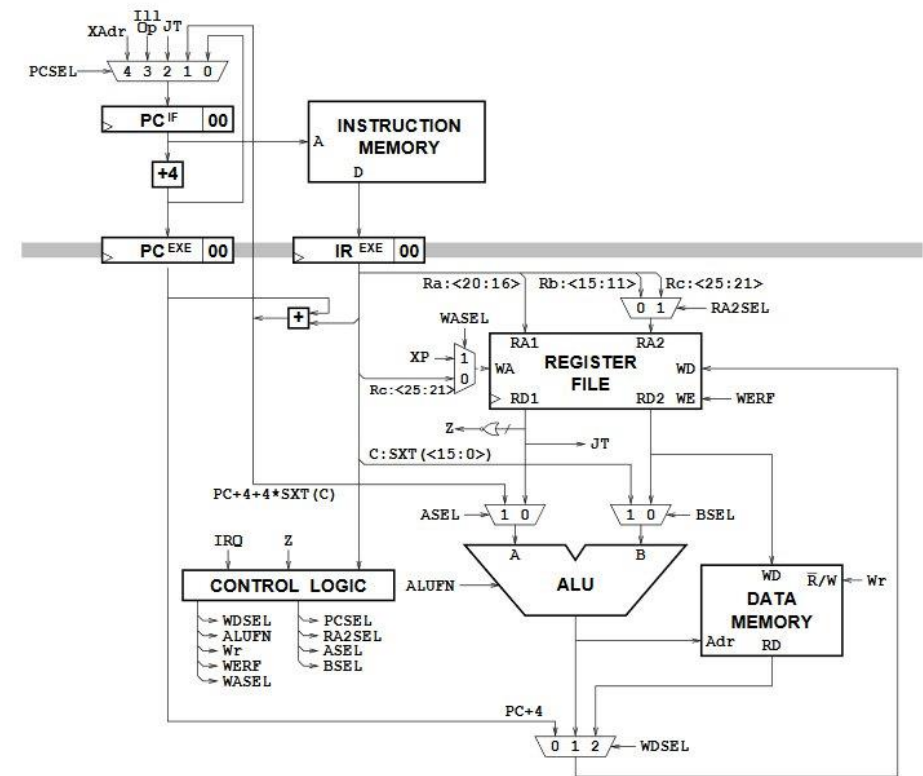
- With pipelines, we are able to speed the machine...
- ... but as we process several instruction at once, some new problems occur
- During this tutorial, we are going to review some of those problems and evaluate their possible solutions

Two-stages pipeline

1. Give a plausible implementation of `NOP()`.
2. Give a *software* solution to the jump problems for each of the following programs:

- | | | | |
|-----|-------------------------------------|-----|------------------------------------|
| (a) | 1 Main: <code>ADDC(R31,0,R1)</code> | (b) | 1 Cas1: <code>ADDC(R1,4,R1)</code> |
| | 2 <code>ADDC(R31,2,R2)</code> | | 2 <code>SUBC(R2,12,R2)</code> |
| | 3 <code>Incr: ADC(R1,5,R1)</code> | | 3 <code>CMPLT(R1,R2,R0)</code> |
| | 4 <code>SUBC(R2,1,R2)</code> | | 4 <code>BNE(R0,Cas2)</code> |
| | 5 <code>BT(R2,Incr)</code> | | 5 <code>MULC(R1,5,R1)</code> |
| | 6 <code>Oper: ADD(R1,R3,R3)</code> | | 6 <code>BR(rtn)</code> |
| | | | 7 Cas2: <code>MULC(R2,5,R2)</code> |
| | | | 8 <code>BR(rtn)</code> |

A β machine with a 2-stage pipeline



Four-stages pipeline

1. Give a software *and* a hardware solution to the data conflicts problems for the following program:

```

1 ADD(R1,R2,R3)
2 SUB(R3,R4,R5)
3 MULC(R2,5,R17)
4 ADD(R5,R1,R1)
5 SUB(R17,R1,R17)
    
```

2. Give a *hardware* solution to the data conflicts problems for the following program:

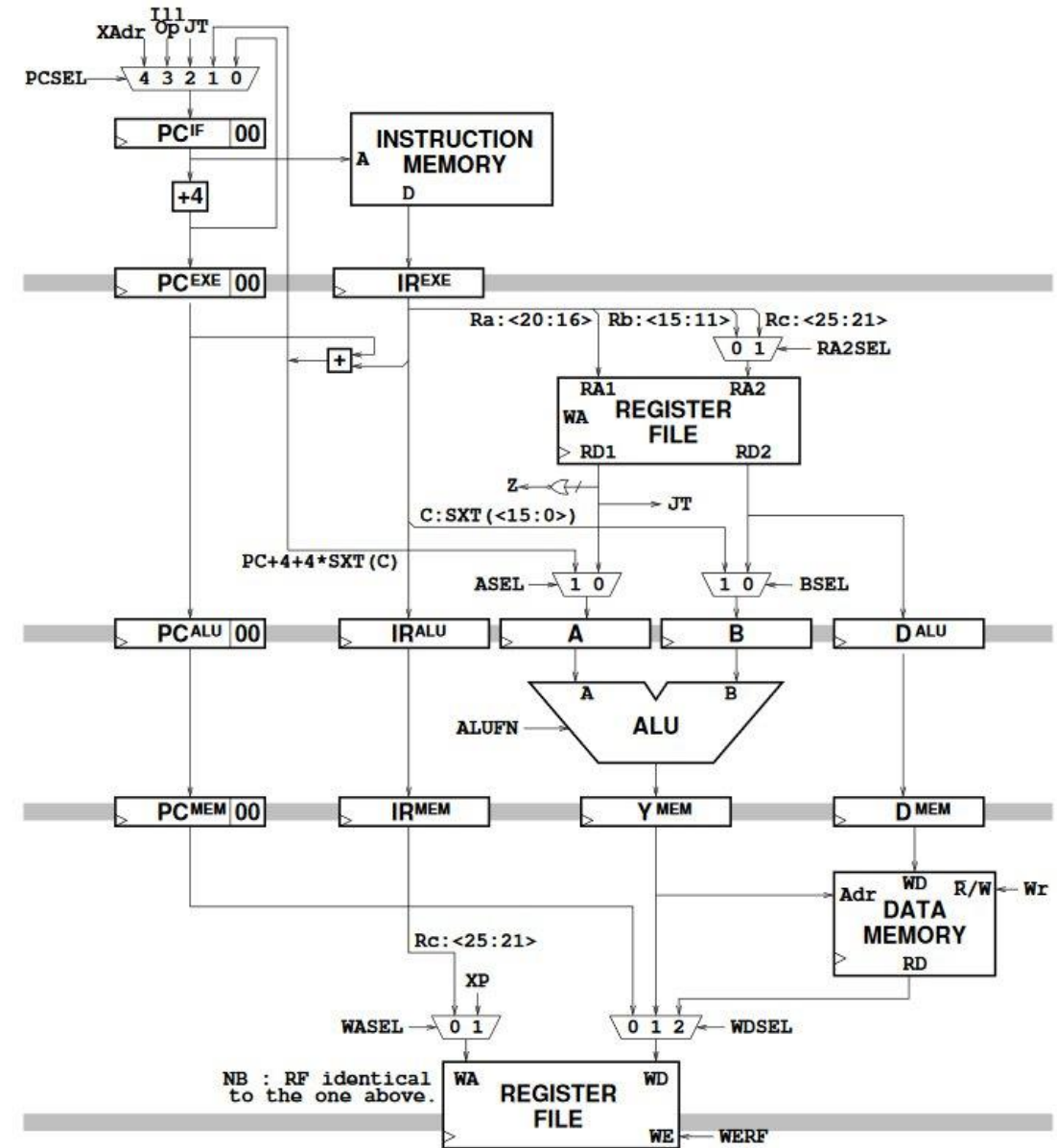
```

1 LD(R1,0,R4)
2 ADD(R1,R4,R5)
3 XOR(R3,R4,R6)
    
```

3. If the β Machine features 2 *bypasses*, what will be the result stored at 0x1000 after the execution of the following program? Why?

```

1 ADDC(R31,3,R0)
2 SUBC(R0,1,R1)
3 MUL(R0,R1,R2)
4 XOR(R0,R2,R3)
5 ST(R3,0x1000,R31)
    
```



Hardware solutions

